Base Station Cooperation for Mission Critical Communications

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Future Wireless Networks Challenges



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Mission Critical Communications



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Mission Critical Communications

- MCC = All communications related to the safety and the security of the civil society. Public safety services, like police forces, firemen, rescue and ambulance services, or employees critical infrastructures, like energy and transportation suppliers
- MCC are conveyed by Professional Mobile Radio (PMR) networks
- One of the most important and indispensable services offered by mission-critical networks is the group communication
- MCC unique requirements : coverage, reliability and secure communications
- Group communication is based on Multimedia Broadcast/Multicast Service (MBMS) in 3GPP

Mission Critical Communications



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Outlines

A Dynamic Clustering Algorithm

- Motivation
- Model
- Problem Formulation
- Algorithms
- Numerical Results

2 Secure Multi-User MIMO Transceiver

- Motivation
- Model
- Transceiver Design
- Numerical Results

3 Conclusions and Future Works

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Conclusions and Future Works

Multicast/Broadcast Single Frequency Network





User SINRs are maximized

Radio resources can be wasted

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Single-Cell Point-To-Multipoint





Degraded user SINRs

Maximized network capacity

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SC-PTM

MBSFN

Degraded users SINRs

Maximized network capacity

Users SINRs are maximized

Radio resources can be wasted

Our objective : design a dynamic clustering algorithm to solve the reliability-capacity tradeoff in MCC

Clustering for cooperative transmissions

- State-of-the-art is mostly on unicast best effort traffic with the goal of maximizing user data rates under static traffic models
- Greedy algorithms based approaches have been widely used in network-centric clustering algorithms [PGH08, YKL16, DdV14, BBB14, Sch17]
- \Rightarrow Not optimal and requires adaptations for group communications
- Dynamic user centric clustering [LHYZ16, GZS14, LZG⁺18, BJIX16, ZWZ⁺17]
- ⇒ Focus on PHY data rate, dynamic traffic constraints ignored, only unicast, no reliability constraint
- MBMS literature [REH11, CCE⁺15]
- ⇒ Do not address the reliability-capacity tradeoff

System Model

- Downlink of a cellular network
- A set $V = \{1, ..., n\}$ of n cells forming an MBSFN synchronization area
- A cluster is a subset S ⊆ V serving a group U of users using multi-point transmissions
- Signal to Interference plus Noise Ratio (SINR) at user $u \in U$:

$$\gamma_u(S) = \frac{\sum_{b \in S} \xi_{ub} g_{ub} P_T}{\sum_{b \in S} (1 - \xi_{ub}) g_{ub} P_T + \sum_{b \notin S} P_T g_{ub} + N_{th}},$$
(1)

where ξ_{ub} denotes the useful portion of the signal received by u from b (see [RHE08] for the detailed calculation)

• For a multicast group \mathcal{U} of N users served by cells in S, the average SINR is :

$$\bar{\gamma}_{\mathcal{U}}(S) = \frac{1}{N} \sum_{u \in \mathcal{U}} \gamma_u(S).$$
⁽²⁾

Dynamic Traffic Model

- Call blocking model : group of users arrive in the synchronization area, use a resource for a group communication for a certain duration and leave the system.
- Poisson arrival (λ), exponential service duration (μ), R resources in every BS.
- When a group arrives in the system, it is served by a subset S of BSs with probability p_S and one resource is consumed in every BS in S.
- We can approximate the blocking probability in BS b by using Erlang-B :

$$\tilde{\Pi}(b) \approx E_B(b, R) = \frac{\frac{\rho_B^A}{R!}}{\sum_{r=0}^R \frac{\rho_b^r}{r!}},$$
(3)

where ρ_b is the load in station *b* (depends on λ , μ and the probability mass function p_s).

Problem Formulation : Cell Clustering

Cell clustering (inner) problem : find a minimizer set S that solves the following set function minimization problem, for a given group U :

$$\min_{S \in \mathcal{P}_{\nu}} \Psi_{\mathcal{U}}(S) \triangleq w(S) - \bar{\gamma}_{\mathcal{U}}(S)$$
(4)

- S : set of serving BSs
- \mathcal{P}_V : set of all subsets of V
- $w(S) = \sum_{b \in S} w_b$: sum of the weights of BSs $b \in S$
- γ
 _{*i*}(S): average SINR of group *U* served by cells in S

 \Rightarrow For a fixed $w = (w_1, \dots, w_n) \in \mathbb{R}^n$, the traffic demand and the clustering policy p_S induces a blocking probability $\tilde{\Pi}(b; w)$ in every b

Problem Formulation : Weights Optimization

Weights optimization (outer) problem : find a minimizer of the quadratic error of the blocking probability wrt target blocking probabilities :

$$\min_{w \in \mathbb{R}^n} G(w) \triangleq \sum_{b=1}^n \|\tilde{\Pi}(b;w) - \bar{\Pi}(b)\|^2$$
(5)

• $\tilde{\Pi}(b; w)$: blocking probability in a BS *b* that depend on the weights vector (*w*)

• $\overline{\Pi}(b)$: target blocking probability that BS *b* should attain

Algorithms

Dynamic Clustering Algorithm (DCA)



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Group Call Cell Clustering Algorithm (GCCA)

$$\min_{S\in\mathcal{P}_{v}}\Psi_{\mathcal{U}}(S)\triangleq w(S)-\bar{\gamma}_{\mathcal{U}}(S)$$

 $\Psi_{\mathcal{U}}$ is a submodular function.

Submodular functions : A set function $F : 2^V \mapsto \mathbb{R}$ is submodular if and only if, for all subsets A, $B \subseteq V$ and $b \in V$ such that $A \subseteq B$ and $b \notin B$, we have : $F(A \cup b) - F(A) \ge F(B \cup b) - F(B)$.





Group Call Cell Clustering Algorithm (GCCA)

$$\min_{S\in\mathcal{P}_{v}}\Psi_{\mathcal{U}}(S)\triangleq w(S)-\bar{\gamma}_{\mathcal{U}}(S)$$

 $\Psi_{\mathcal{U}}$ is a submodular function.

- Minimizing $\Psi_{\mathcal{U}}$ is a **submodular minimization** problem.
- This problem is solved using the Minimum-Norm Algorithm [Fuj84, Wol76].
- The performance of this solution is compared to a Greedy based approach.

Cell Weights Optimization Algorithm (CWOA)

$$\min_{w\in\mathbb{R}^n}G(w)\triangleq\sum_{b=1}^n\|\tilde{\Pi}(b;w)-\bar{\Pi}(b)\|^2$$

- We rely on direct search methods to minimize G.
- These methods provide simpler calculations and relatively low storage requirements over derivative based methods.
- The most popular is the Nelder-Mead simplex method.

Cell Weights Optimization Algorithm (CWOA)

$$\min_{w \in \mathbb{R}^n} G(w) \triangleq \sum_{b=1}^n \|\tilde{\Pi}(b;w) - \bar{\Pi}(b)\|^2$$

- Nelder-Mead algorithm can stagnate, fail to converge or converge to a non-optimal vertex.
- A possible improvement of the original algorithm is to impose restarts of the algorithm during the optimization run.
- An oriented restart of the Nelder-Mead algorithm adapted to our model :
 - If $\Pi(b; w)$ is too small, weight w_b is decreased randomly;
 - If $\Pi(b; w) > \Pi(b)$, weight w_b is increased randomly.

Numerical Results : Group Call Clustering



Figure – Evolution of $\Psi_{\mathcal{U}}$ along the iterations of the minimum-norm algorithm

SC-PTM and full MBSFN cooperation schemes are outperformed by the proposed algorithm in very few iterations.

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Numerical Results : Cell Weights Optimization



Evolution of ${\it G}$ and cell blocking probabilities along the iterations of the Nelder-Mead algorithm

The reliability-capacity tradeoff is well handled by the proposed scheme.

Numerical Results : SINR Improvements



Users and group mean SINR CDF

The proposed scheme lies in-between MBSFN and SC-PTM in terms of SINR.

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Conclusion

 \Rightarrow Our algorithm is able to adapt to traffic variations by maximizing the coverage under the constraint of a blocking probability.

Related publications :

- A. Daher, M. Coupechoux, P. Godlewski, P. Ngouat, P. Minot, A Dynamic Clustering Algorithm for Multi-Point Transmissions in Mission-Critical Communications in Mission-Critical Communications, IEEE Trans. on Wireless Communications, Apr. 2020.
- Alaa Daher, M. Shabbir Ali, Marceau Coupechoux, Philippe Godlewski, Pierre Ngouat, Pierre Minot, A Repetition Scheme for MBSFN Based Mission-Critical Communications, IEEE VTC Fall, Aug. 2018.
- A. Daher, M. Coupechoux, P. Godlewski, J.M. Kélif, P. Ngouat, P. Minot, SINR Model for MBSFN Based Mission Critical Communications, IEEE VTC Fall, Sept. 2017.
- A. Daher, M. Coupechoux, P. Godlewski, P. Ngouat, P. Minot, SC-PTM or MBSFN for Mission Critical Communications?, IEEE VTC Spring, June 2017.

Image: Image:

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- The previous study is ignoring security aspects and relies on simple physical layer models.
- ⇒ Our goal now : design a multi-BS multi-antenna transceiver for MCC that is robust to CSI errors (*reliability requirement of MCC*) and that is secured with respect to the presence of multiple eavesdroppers (*security*).
 - Specifically, we incorporate security in two ways :
 - MIMO beamforming is used to achieve the desired performance gain at legitimate users while degrading eavesdroppers channel;
 - Artificial Noise (AN) is added at the transmitter to guarantee additional security over the designed transceivers.
 - Robustness is considered wrt :
 - Stochastic errors
 - Norm-bounded errors

Network Model



Figure – Blue cells are serving a group of legitimate users (green diamonds); multiple eavesdroppers may overhear the communication (red stars).

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Network Model

A greedy clustering is adopted for simplicity :

Algorithm 1 Greedy Clustering

1: Input : Locations of BSs and group users, $K'_T \leq |\mathcal{B}|$: minimum cluster size

- 2: Init : $\mathcal{S} \leftarrow \emptyset$
- 3: for every user do
- 4: Find the BS t providing the highest receive power
- 5: $\mathcal{S} \leftarrow \mathcal{S} \cup \{t\}$
- 6: end for
- 7: if $|\mathcal{S}| < K'_T$ then
- 8: Find the set S' of $K'_T |S|$ BSs maximizing the sum SINR for group users
- 9: $\mathcal{S} \leftarrow \mathcal{S} \cup \mathcal{S}'$
- 10: end if
- 11: return S

Transceiver Model



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Transceiver Model

• The signal transmitted from *t*-th BS is given by :

$$\mathbf{x}_t = \mathbf{V}_t \mathbf{d} + \mathbf{W}_t \mathbf{z}_t, \tag{6}$$

where \mathbf{V}_t is a precoder, \mathbf{z}_t is an zero-mean additional AN vector with variance $\mathbb{E}[\mathbf{z}_t \mathbf{z}_t^H] = \sigma_{zt}^2 \mathbf{I}_{N_T}$ and \mathbf{W}_t is an AN-shaping matrix.

• The estimated data at user I is :

$$\widehat{\mathbf{d}}_{l} = \mathbf{R}_{l} \mathbf{C}_{tl} \mathbf{V}_{t} \mathbf{d} + \mathbf{R}_{l} \mathbf{C}_{tl} \mathbf{W}_{t} \mathbf{z}_{t} + \mathbf{R}_{l} \mathbf{n}_{l}.$$
(7)

where \mathbf{R}_{l} is a receive filter.

The MSE at the *I*-th legitimate UE is :

$$\epsilon_I \triangleq \mathbb{E}[||\mathbf{d} - \widehat{\mathbf{d}}_I||^2], \tag{8}$$

In the same way, the MSE at eavesdropper e is :

$$\epsilon_e \triangleq \mathbb{E}[||\mathbf{d} - \overline{\mathbf{d}}_e||^2]. \tag{9}$$

Transceiver Model

CSI errors :

- At legitimate users : $\mathbf{G}_{te} = \widehat{\mathbf{G}}_{te} + \mathbf{\Delta}_{te}$
- At eavesdroppers : $C_{tl} = \widehat{C}_{tl} + \Delta_{tl}$

Assumption (Stochastic error model)

CSI errors Δ_{te} and Δ_{tl} are modeled as Gaussian random variables such that $\mathbb{E}[\Delta_{te}\Delta_{te}^{H}] = \sigma_{te}^{2}I$ and $\mathbb{E}[\Delta_{tl}\Delta_{tl}^{H}] = \sigma_{tl}^{2}I$.

Assumption (Norm-bounded error model)

CSI errors Δ_{te} and Δ_{tl} are respectively taken in continuous sets, called uncertainty regions :

$$\mathcal{G}_{te} = \{ \boldsymbol{\Delta}_{te} : ||\boldsymbol{\Delta}_{te}||^2 \le \tau_{te} \},$$
(10)

$$\mathcal{C}_{tl} = \{ \boldsymbol{\Delta}_{tl} : ||\boldsymbol{\Delta}_{tl}||^2 \le \tau_{tl} \},$$
(11)

where τ_{te} and τ_{tl} denote the radii of the uncertainty regions.

Stochastic Errors : Problem Formulation

In presence of stochastic errors :

$$\begin{array}{ll} \underset{t=1...K_{T}, l=1...K_{R}}{\underset{t=1}{\text{minimize}}} & \sum_{l=1}^{K_{R}} \epsilon_{l}, \\ \text{subject to} & C1: \ \epsilon_{e} \geq \Gamma, \quad \forall e \in \{1, \cdots, K_{E}\}, \\ & C2: \ P_{t} \leq P_{T}, \quad \forall t \in \{1, \cdots, K_{T}\}, \end{array}$$

$$(\mathcal{P}_{1})$$

where Γ is a design parameter that represents a lower bound on the achievable MSE expected at each eavesdropper.

Stochastic Errors : Stationary Point

Proposition

$$\mathbf{V}_{t} = (\mathbf{A}_{t})^{-1} \Big(\sum_{l=1}^{K_{R}} \widehat{\mathbf{C}}_{tl}^{H} \mathbf{R}_{l}^{H} - \sum_{e=1}^{K_{E}} \lambda_{e} \widehat{\mathbf{G}}_{te}^{H} \mathbf{E}_{e}^{H} \Big)$$
(12)

$$\mathbf{W}_{t} = \mathbf{B}_{t} / \sqrt{[\mathrm{tr}(\mathbf{B}_{t}\mathbf{B}_{t}^{\mathrm{H}})]}$$
(13)

$$\mathbf{R}_{l} = \left(\sum_{t=1}^{K_{T}} \mathbf{V}_{t}^{H} \widehat{\mathbf{C}}_{tl}^{H}\right) \left(\sum_{t=1}^{K_{T}} \widehat{\mathbf{C}}_{tl} \mathbf{V}_{t} \mathbf{V}_{t}^{H} \widehat{\mathbf{C}}_{tl}^{H} + \sum_{t=1}^{K_{T}} \sigma_{zt}^{2} \widehat{\mathbf{C}}_{tl} \mathbf{W}_{t} \mathbf{W}_{t}^{H} \widehat{\mathbf{C}}_{tl}^{H} + \sum_{t=1}^{K_{T}} \sigma_{zt}^{2} \widehat{\mathbf{C}}_{tl} \mathbf{W}_{t}$$

$$\sigma_{nl}^{2}\mathbf{I} + \sum_{t=1}^{N_{f}} \sigma_{tl}^{2} \operatorname{tr}(\mathbf{V}_{t}\mathbf{V}_{t}^{H})\mathbf{I} + \sum_{t=1}^{N_{f}} \sigma_{tl}^{2} \sigma_{zt}^{2} \operatorname{tr}(\mathbf{W}_{t}\mathbf{W}_{t}^{H})\mathbf{I} \Big)^{-1}$$
(14)

$$\mathbf{B}_{t} = \mathbf{I} - \mathbf{A}_{t}^{H} (\mathbf{A}_{t} \mathbf{A}_{t}^{H})^{-1} \mathbf{A}_{t},$$

$$\mathbf{A}_{t} = \sum_{l=1}^{K_{R}} \widehat{\mathbf{C}}_{tl}^{H} \mathbf{R}_{l}^{H} \mathbf{R}_{l} \widehat{\mathbf{C}}_{tl} + \sum_{l=1}^{K_{R}} \sigma_{tl}^{2} \operatorname{tr}(\mathbf{R}_{l} \mathbf{R}_{l}^{H}) - \sum_{e=1}^{K_{E}} \lambda_{e} \widehat{\mathbf{G}}_{te}^{H} \mathbf{E}_{e} \widehat{\mathbf{G}}_{te} - \sum_{e=1}^{K_{E}} \lambda_{e} \sigma_{te}^{2} \operatorname{tr}(\mathbf{E}_{e} \mathbf{E}_{e}^{H}) + \lambda_{t}^{'} \mathbf{I}.$$

Stochastic Errors : Block Descent Algorithm

Algorithm 2 Block Descent Algorithm

- 1: Input : β , K_T , K_R , K_E , $\widehat{\mathbf{C}}_{tl}$, $\widehat{\mathbf{G}}_{te}$, σ_{nl} , σ_{ne} , $P_t \forall t \in \{1, \dots, K_T\}$, $l \in \{1, \dots, K_R\}$, and $e \in \{1, \dots, K_E\}$
- 2: Init : Randomly generate V_t , $W_t \forall t \in \{1, \cdots, K_T\}$, $\epsilon'_l \leftarrow 0$, $\epsilon_l \leftarrow 0$, $\forall l \in \{1, \cdots, K_R\}$

3: repeat

4:
$$\epsilon'_I \leftarrow \epsilon_I, \forall I \in \{1, \cdots, K_R\}$$

- 5: Update $\mathbf{E}_e \ \forall e \in \{1, \cdots, K_E\}$ (MMSE filter)
- 6: Update \mathbf{R}_l using \mathbf{V}_t , \mathbf{W}_t in (14) $\forall l \in \{1, \cdots, K_R\}$
- 7: Solve for λ_e and $\lambda_t^{'}$ using C1, C2 $\forall t \in \{1, \cdots, K_T\}$, and $\forall e \in \{1, \cdots, K_E\}$
- 8: Update \mathbf{V}_t using λ_e , λ'_t , \mathbf{R}_l , \mathbf{E}_e in (12) $\forall t = \{1, \cdots, K_T\}$
- 9: Update \mathbf{W}_t using \mathbf{V}_t in (13) $\forall t = \{1, \cdots, K_T\}$
- 10: Compute ϵ_l using \mathbf{V}_t , λ_e , λ'_t , \mathbf{W}_t , and \mathbf{R}_l

11: until
$$|\epsilon_l - \epsilon'_l| \leq \beta$$
, $\forall l \in \{1, \cdots, K_R\}$

12: return V_t , W_t , $\forall t \in \{1, \cdots, K_T\}$, R_l , $\epsilon_l \ \forall l \in \{1, \cdots, K_R\}$

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Norm-Bounded Errors : Problem Formulation

In presence of norm-bounded errors :

$$\begin{array}{ll} \underset{\substack{V_{t}, W_{t}, R_{l} \\ t=1...K_{T}, l=1...K_{R} \end{array}}{\text{minimize}} & \sum_{l=1}^{K_{R}} \epsilon_{l}, \\ \text{subject to} & C1, C2 & (\mathcal{P}_{2}) \\ & C4: \ \mathbf{\Delta}_{te} \in \mathcal{G}_{te}, \quad \forall t \in \{1, \cdots, K_{T}\}, \forall e \in \{1, \cdots, K_{E}\} \\ & C5: \ \mathbf{\Delta}_{tl} \in \mathcal{C}_{tl}, \quad \forall t \in \{1, \cdots, K_{T}\}, \forall l \in \{1, \cdots, K_{R}\} \end{array}$$

An equivalent (robust) formulation :

$$\begin{array}{ll} \underset{\substack{v_{t}, w_{t}, R_{l} \\ t=1...K_{T}, l=1...K_{R}}{\text{min}} & \sum_{l=1}^{K_{R}} \epsilon_{l}, \\ \text{subject to} & C2, \\ & C6: & \sum_{\boldsymbol{\Delta}_{te} \in \mathcal{G}_{te}} \epsilon_{e} \geq \Gamma, \quad \forall e \in \{1, \cdots, K_{E}\} \end{array}$$

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Norm-Bounded Errors : Problem Decomposition

Sub-problem $\bar{\mathcal{P}}'_2$: Assume that the worst-case channel errors Δ_{te} and Δ_{tl} are known and solve for the transceiver matrices :

$$\begin{array}{ll} \underset{t=1...K_{T}, l=1...K_{R}}{\text{minimize}} & \sum_{l=1}^{K_{R}} \epsilon_{l}, \\ \underset{t=1...K_{T}, l=1...K_{R}}{\text{minimize}} & (\bar{\mathcal{P}}_{2}') \end{array}$$
subject to $C1, C2$

Sub-problem $\bar{\mathcal{P}}_2''$: Assume that transceiver matrices and the worst-case error Δ_{te} are known and solve for the worst-case error Δ_{tl} .

$$\begin{array}{ll} \underset{t=1...K_{T}, l=1...K_{R}}{\text{minimize}} & -\sum_{l=1}^{K_{R}} \epsilon_{l}, \\ \text{subject to} & C5: ||\boldsymbol{\Delta}_{tl}||^{2} \leq \tau_{tl} & \forall t \in \{1, \cdots, K_{T}\}, \forall l \in \{1, \cdots, K_{R}\} \end{array}$$

Norm-Bounded Errors : Problem Decomposition

Sub-problem $\bar{\mathcal{P}}_{2}^{\prime\prime\prime}$: Assume that transceiver matrices and the worst-case error $\Delta_{t/}$ are known and solve for the worst-case error Δ_{te} .

$$\begin{array}{ll} \underset{\boldsymbol{\Delta}_{te},t=1...K_{T}}{\text{minimize}} & \epsilon_{e} \geq \Gamma, \\ \text{subject to} & C4: ||\boldsymbol{\Delta}_{te}||^{2} \leq \tau_{te} & \forall t \in \{1,\cdots,K_{T}\} \end{array}$$

Norm-Bounded Errors : Algorithm

Algorithm 3 Iterative Algorithms for NBE

- 1: Input : β , K_T , K_R , K_E , \mathbf{E}_e , $\widehat{\mathbf{C}}_{tl}$, $\widehat{\mathbf{G}}_{te}$, τ_{tl} , τ_{te} , P_t , $\forall t \in \{1, \dots, K_T\}$, $l \in \{1, \dots, K_R\}$, and $e \in \{1, \dots, K_E\}$
- 2: Init : Randomly generate V_t , W_t , $\forall t \in \{1, \dots, K_T\}$, $\Delta_{t'} \in C_{t'}$, $\Delta_{te} \in G_{te}$, $\epsilon_l \leftarrow 0$, $\forall t \in \{1, \dots, K_T\}$, $l \in \{1, \dots, K_R\}$, $e \in \{1, \dots, K_E\}$,
- 3: repeat
- 4: $\epsilon'_{I} \leftarrow \epsilon_{I}, \forall I \in \{1, \cdots, K_{R}\}$
- 5: Solve \overline{P}'_2 and update V_t , W_t , R_l using Δ_{tl} , Δ_{te} , E_e and Algorithm 2 $\forall t \in \{1, \cdots, K_T\}$, $l \in \{1, \cdots, K_R\}$, $e \in \{1, \cdots, K_E\}$
- 6: Solve $\overline{\mathcal{P}}_{2}^{\prime\prime}$ and update Δ_{tl} using V_{t} , W_{t} , $R_{l} \forall t \in \{1, \cdots, K_{T}\}$, $l \in \{1, \cdots, K_{R}\}$
- 7: Solve $\bar{\mathcal{P}}_{2}^{\prime\prime\prime}$ and update Δ_{te} using V_t , W_t , $E_e \ \forall t \in \{1, \cdots, K_T\}$, $e \in \{1, \cdots, K_E\}$
- 8: Compute ϵ_l using V_t , W_t , and R_l , Δ_{tl} , Δ_{te}

9: until $|\epsilon_l - \epsilon'_l| \leq \beta$, $\forall l \in \{1, \cdots, K_R\}$

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Numerical Results : Physical Layer



- Robust design improves the reliability of MCC.
- Due to increased uncertainty, norm-bounded errors leads to poorer performances compared to stochastic errors.

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Numerical Results : Physical Layer



• System performance decreases with increasing CSI error variance.

The higher the variance, the higher the gap bw robust and non-robust.

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Numerical Results : Physical Layer



• Security gap :
$$S_g = SNR_{min}^L - SNR_{max}^E$$

- Below *SNR^E_{max}* at eavesdroppers, the communication is secure, above *SNR^L_{min}* at legitimate UEs, the communication is reliable.
- NR : 4dB, R-NBE : 2dB, R-SE : -2dB at BER 10^{-4} for leg. UEs and 0.3 for eaves
- Robust designs achieve reduced security gap (enhanced secrecy performance).

Physical Level Simulation Results



• Presence of AN degrades the performance at eavesdropper whereas it has a very low impact on BER of legitimate UE.

Numerical Results : System Level



- All the systems, MBSFN, SC-PTM and dynamic clustering demonstrate enhanced performance for legitimate UEs as compared to eavesdroppers.
- At legitimate UEs, MBSFN performs the best, SC-PTM performs worst, whereas dynamic clustering shows a good trade off.

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Conclusion

 $\Rightarrow\,$ We designed an efficient multi-BS multi-antenna secure MIMO transceiver that meets the unique requirements of MCC.

Related publications :

- Deepa Jagyasi, Marceau Coupechoux, Secure Transceiver Design for Multi-user MIMO Multicast Mission Critical Communication System, *submitted*.
- Deepa Jiagyasi, Alaa Daher, Marceau Coupechoux, Multicell MIMO Transceiver Design for Mission-Critical Communication, IEEE Globecom, Dec. 2019.

Outline

1 A Dynamic Clustering Algorithm

- Motivation
- Model
- Problem Formulation
- Algorithms
- Numerical Results

2 Secure Multi-User MIMO Transceiver

- Motivation
- Model
- Transceiver Design
- Numerical Results

3 Conclusions and Future Works

Conclusions and Future Works

- Mission critical communications for public safety, critical infrastructure, etc is a relatively unexplored topic.
- They deserve specific designs because of a unique set of joint constraints like security, reliability, coverage, multicast services.
- Traditional designs focusing on capacity and user data rate maximization no longer apply.
- Future works includes mm waves studies and new approaches to reduce energy consumption.

Thank you for your attention !

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