

# Optimal Trajectories of a UAV Base Station Using Lagrangian Mechanics

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**Abstract**—This paper considers the problem of optimizing the trajectory of an Unmanned Aerial Vehicle (UAV) Base Station (BS). A map is considered, characterized by a traffic intensity of users to be served. The UAV BS must travel from a given initial location at an initial time to a final position within a given duration and serves the traffic on its way. The problem consists in finding the optimal trajectory that minimizes a certain cost depending on the velocity and on the amount of served traffic. The problem is formulated using the framework of Lagrangian mechanics. When the traffic intensity is quadratic (single-phase), we derive closed-form formulas for the optimal trajectory. When the traffic intensity is bi-phase, necessary conditions of optimality are provided and an Alternating Optimization Algorithm is proposed, that returns a trajectory satisfying these conditions. The Algorithm is initialized with a Model Predictive Control (MPC) online algorithm. Numerical results show how the trajectory is improved with respect to the MPC solution.

## I. INTRODUCTION

Unmanned Aerial Vehicles (UAV) are expected to play an increasing role in future wireless networks<sup>1</sup> [1]. UAVs may be deployed in an ad hoc manner when the traditional cellular infrastructure is missing. They can serve as relays to reach distant users outside the coverage of wireless networks. They also may be used to disseminate data to ground stations or collect information from sensors. In this paper, we address one of the envisioned use cases for UAV-aided wireless communications, which relates to cellular network offloading in highly crowded areas [1]. More specifically, we focus on the path planning problem or trajectory optimization problem that consists in finding an optimal path for a UAV Base Station (BS) that minimizes a certain cost depending on the velocity and on the amount of served traffic. Our approach is based on the Lagrangian mechanics framework.

### A. Related Work

UAV trajectory optimization for networks has been tackled maybe for the first time in [2]. The model consists in a UAV flying over a sensor network from which it has to collect some data. The problem consists in optimizing the trajectory length of the UAV under the constraint that it collects the required amount of data from every sensor. Authors use a reinforcement learning approach where improved trajectories are sequentially learned over several tour iterations. This model is different from ours as it allows the UAV to learn the optimal trajectory from previous experience. The problem of optimally deploying UAV BSs to serve traffic

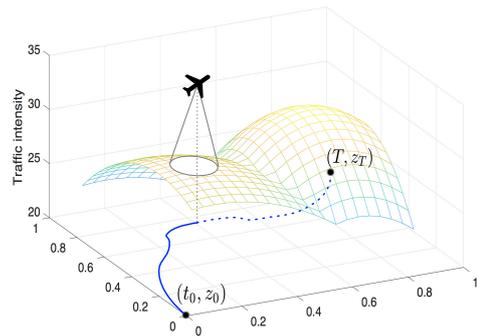


Fig. 1: A UAV Base Station travels from  $z_0$  at  $t_0$  to  $z_T$  at  $T$  and serves a user traffic characterized by its intensity.

demand has been addressed in the literature by considering static UAVs BSs or relays, see e.g. [3], [4]. The goal is to optimally position the UAV so as to maximize the data rate with ground stations or the number of served users. In a very recent work [5], a data rate-energy trade-off is studied. In these works the notion of trajectory is either ignored or restricted to be circular or linear. In robotics and autonomous systems, trajectory optimization is known as *path planning* [6]. In this aim, there are classical methods like Cell Decomposition, Potential Field Method or Probabilistic Road Map and there are heuristic approaches, e.g. bio-inspired algorithms. Authors of [7] have capitalized on this literature and proposed a path planning algorithm for drone BSs based on A\* algorithm. The main goal of these papers is to reach a destination while avoiding obstacles, and in [7] the speed cannot be controlled. In our work, we intend to minimize a certain cost function along the trajectory by controlling the velocity of the UAV. This goal is studied in optimal control theory [8] and is applied for example in aircraft trajectory planning [9]. Most numerical methods in control theory can be classified in *direct* and *indirect* methods. Indirect methods provide analytical solutions from the calculus of variations and use first order necessary conditions for a trajectory to be optimal. In direct methods, the problem is transformed in a non linear programming problem using discretized time, locations and controls. Direct methods are heavily applied in a series of very recent publications in the field of UAV-aided communications. In [10] for example, a UAV relay assists the communication between a source and a destination. As

<sup>1</sup>J. Darbon is supported by NSF DMS-1820821.

the resulting problem is non-convex, it is first approximated and then solved by successive convex optimization. In [11], the objective is to maximize the energy efficiency of a UAV-to-ground station communication by taking into account the propulsion energy consumption and by optimizing the trajectory. Again, sequential convex optimization is applied to an approximated problem. In the same vein, [12] considers multiple-UAV BSs used to serve fixed users. The quality of the solution to the nonlinear program may heavily depend on the initial guess. Authors thus propose an heuristic based on circular trajectories to initialize their algorithm. With direct methods, because of the discretization, the differential equations and the constraints of the systems are satisfied only at discrete points. This can lead to less accurate solutions than indirect methods and the quality of the solution depends on the quantization step [13]. Although every iteration of the sequential convex optimization technique has a polynomial time complexity, practical resolution time may dramatically increase with the quantization grid and the dimension of the problem. We thus propose in this paper an indirect approach based on Lagrangian mechanics that has the advantage to provide closed-form expressions for the optimal trajectories when the potential is quadratic (we say *single-phase*). When the potential is quadratic by region (or *multi-phase*) the optimization process consists in finding the right crossing time and location on the interface of the regions. This question is an active field of research in control theory, see e.g. [14]. As explained in [15], [16], a trajectory optimization problem can be decomposed in different *phases* or *arcs*. Phases are sequential in time, i.e., they partition the time domain. Differential equations describing the system dynamics cannot change during a phase. This point of view allows us to consider the multi-phase problem.

## B. Contributions

Our contributions are the following:

- *Problem Formulation:* To the best of our knowledge, this is the first time that the UAV BS trajectory problem is formulated using the formalism of Lagrangian mechanics. This approach provides closed-form equations when the potential is quadratic and thus very low complexity solutions compared to existing solutions in the literature.
- *Closed-form expression of the optimal trajectory with single phase traffic intensity:* When the traffic intensity map is made of a single hot spot or traffic hole, has a quadratic form (*single phase*), and is time-independent, closed form expressions for the optimal trajectory are derived. It consists in a part of hyperbole for a hot spot and corresponds to a repulsor in mechanics. For a traffic hole, the trajectory is on an ellipse and corresponds to the case of an attractor in mechanics.
- *Characterization of the optimal solution in multi-phase traffic intensity:* When the traffic map has several hot spots or traffic holes (*multi-phase*) whose regions are separated by interfaces and is time-independent, we derive necessary conditions to be fulfilled by the position and the instant at which the optimal trajectory crosses an interface (see Theorem 2).

- *An online algorithm for multi-phase time-varying traffic intensity:* When the traffic map is multi-phase and is time-varying, we propose an online algorithm based on MPC.
- *An Alternating Optimization Algorithm for bi-phase time-independent traffic intensity:* When the traffic intensity is made of two hot spots separated by an interface (*bi-phase*) and is time-independent, we propose an Alternating Optimization Algorithm that finds a stationary point for the cost function. This algorithm has a complexity  $O(1)$  at every iteration, whereas iterations of the sequential convex optimization technique have polynomial time complexity (see Algorithm 1).

The paper is structured as follows. In Section II we give the system model and its interpretation in terms of Lagrangian mechanics. In Section III, we formulate the problem and give preliminary results. Section IV is devoted to the characterization of the optimal trajectories. Section V presents our algorithms and Section VI concludes the paper. All proofs are given in the Appendix of [17].

**Notations:** Let  $f : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$  defined by  $f(x, y)$  where  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$  and  $y = (y_1, \dots, y_m) \in \mathbb{R}^m$ . Let  $a \in \mathbb{R}^n$  and  $b \in \mathbb{R}^m$ . We denote by  $\frac{\partial f}{\partial x_i}(a, b)$  the partial derivative of  $f$  with respect to the variable  $x_i$  at  $(a, b) \in \mathbb{R}^n \times \mathbb{R}^m$ . We also introduce the notations  $\nabla_x f(a, b) = (\frac{\partial f}{\partial x_1}(a, b), \dots, \frac{\partial f}{\partial x_n}(a, b)) \in \mathbb{R}^n$  and  $\nabla_y f(a, b) = (\frac{\partial f}{\partial y_1}(a, b), \dots, \frac{\partial f}{\partial y_m}(a, b)) \in \mathbb{R}^m$ .

## II. SYSTEM MODEL AND INTERPRETATION

### A. System Model

We consider a network area characterized by a traffic density at position  $z$  and time  $t$ . We intend to control the trajectory and the velocity of a UAV base station, which is located in  $z_0 \triangleq z(t_0)$  at  $t_0$  and shall reach a destination  $z_T \triangleq z(T)$  at  $T$  with the aim of minimizing a cost determined by the velocity and the traffic, defined hereafter by (1). At  $(t, z)$ , we assume that the UAV BS is able to cover an area, from which it can serve users (see Figure 1). The velocity of the UAV BS induces an energy cost. In this model, we control the velocity vector  $a$  of the UAV BS. The general form of the cost function is as follows

$$\mathcal{L}(t, z, a) = \frac{K}{2} \|a\|^2 - u(t, z) \quad (1)$$

where the first term is a cost related to the velocity of the vehicle ( $K$  is a positive constant), and  $\|\cdot\|$  denotes the usual Euclidean norm. The higher is the speed, the higher is the energy cost. The second term is a *user traffic intensity*, i.e., the amount of traffic served by the UAV BS at  $(t, z)$ . Note that a non-zero energy at null speed can be incorporated in the model by adding a constant. Without loss of generality, we assume that this constant is null.

## III. LAGRANGIAN MECHANICS FORMULATION

### A. Problem Formulation

Let  $S(t_0, z_0, T, z_T)$  be the minimal total cost along any trajectory between  $z_0$  at  $t_0$  and  $z_T$  at  $T$  (also called *the action* in mechanics or *value function* in control theory).

Let us define  $\Omega(t_0, T)$  as the space of absolutely continuous functions from  $[t_0; T]$  to  $\mathbb{R}^2$ . Our problem can now be formulated as follows

$$S(t_0, z_0, T, z_T) = \min_{a \in \Omega(t_0, T)} \int_{t_0}^T \mathcal{L}(s, z(s), a(s)) ds + J(z(T)) \quad (2)$$

where  $\frac{dz}{dt}(t) = a(t)$ ,  $z(t_0) = z_0$ , and  $J$  is the terminal cost defined by  $J(z) = 0$  if  $z = z_T$  and  $J(z) = +\infty$  otherwise. For simplicity reasons, we assume the existence and uniqueness of the optimal control  $a^*(t)$  in (2) and denote the associated optimal trajectory  $z^*(t)$ . In a traffic map symmetric with respect to  $z_0$  and  $z_T$ , the reader can convince himself that the uniqueness is not guaranteed.

### B. Preliminary Results From Lagrangian Mechanics

We provide in this section important results from the Lagrangian mechanics for the convenience of the reader.

**Definition 1** (Impulsion). *The impulsion function is defined as*

$$p(t, z, a) := \nabla_a \mathcal{L}(t, z, a) \quad (3)$$

In the Newtonian classical framework that is used here (see (1)), the impulsion is the product of the particle mass by its velocity (hence the standard term "impulsion").

**Definition 2.** *The Hamiltonian function is defined as*

$$H(t, z, p) := \max_{a \in \mathbb{R}^2} p \cdot a - \mathcal{L}(t, z, a). \quad (4)$$

**Lemma 1** (Euler-Lagrange Equations). *Along the optimal trajectory  $z^*(t)$  that starts from  $z_0$  at  $t_0$  and ends at  $z_T$  at  $T$ , we have*

$$\frac{d}{dt} \nabla_a \mathcal{L}(t, z^*(t), a^*(t)) = \nabla_z \mathcal{L}(t, z^*(t), a^*(t)) \quad (5)$$

or equivalently

$$\frac{dp}{dt}(t, z^*(t), a^*(t)) = \nabla_z \mathcal{L}(t, z^*(t), a^*(t)) \quad (6)$$

The Euler-Lagrange equation is the first-order necessary condition for optimality and holds for every point on the optimal trajectory.

**Lemma 2.** *If the Lagrangian  $\mathcal{L}(t, z, a)$  is time-independent and  $\alpha$ -homogeneous in  $z$  and  $a$  for  $\alpha > 0$ , i.e.,  $\mathcal{L}(\lambda z, \lambda a) = |\lambda|^\alpha \mathcal{L}(z, a)$  for all  $\lambda \in \mathbb{R}$ ,  $S$  given by (2) reads*

$$S(t_0, z_0, T, z_T) = \frac{1}{\alpha} [z \cdot p]_{t_0}^T + J(z_T). \quad (7)$$

**Lemma 3** (Hamilton-Jacobi). *Along the optimal trajectory, we have for  $t \in (t_0; T)$*

$$\frac{\partial S}{\partial t_0}(t, z^*(t), T, z_T) = H(t, z^*(t), -p^*(t)) \quad (8)$$

$$\frac{\partial S}{\partial T}(t_0, z_0, t, z^*(t)) = -H(t, z^*(t), p^*(t)) \quad (9)$$

where

$$p^*(t) = \nabla_a \mathcal{L}(t, z^*(t), a^*(t)) = \nabla_z S(t, z^*(t), T, z_T) \quad (10)$$

From now, we assume that the Lagrangian is time-independent, i.e.,  $\mathcal{L}(t, z, a) = \mathcal{L}(z, a)$ , and is an even function in  $a$ , i.e.,  $\mathcal{L}(z, -a) = \mathcal{L}(z, a)$ . A direct consequence is that  $H$  is time-independent and is an even function in  $p$ , i.e., we write  $H(t, z, p) = H(z, p)$  and  $H(z, -p) = H(z, p)$ .

## IV. OPTIMAL TRAJECTORY

In this section, we characterize the optimal trajectory when the traffic intensity is a quadratic form and also when it is made of two regions of quadratic form separated by an interface<sup>2</sup>. We call these two cases *single-phase* and *multiple-phase* intensities respectively. Both cases satisfy our assumptions on the Lagrangian with  $\alpha = 2$ .

### A. Single-Phase Optimal Trajectory

Assume that the traffic intensity is of the form  $u(z) = \frac{1}{2}u_0||z||^2$ . When  $u_0 > 0$ , this function models a traffic hole in  $z = 0$ . When  $u_0 < 0$ , it models a traffic hot spot at  $z = 0$ . We disregard the case  $u_0 = 0$  because it corresponds to a constant traffic intensity that is not of interest in this paper. Thus the cost function has the following form

$$\mathcal{L}(z, a) = \frac{1}{2}K||a||^2 - \frac{1}{2}u_0||z||^2 \quad (11)$$

Note that

$$p(z, a) = \nabla_a \mathcal{L}(z, a) = Ka \quad (12)$$

1) *Trajectory Equation:* In the single phase case, we have a closed form expression of the trajectory.

**Theorem 1.** *If  $u_0 < 0$ , the cost function is given by (13), the optimal trajectory is*

$$z^*(t) = \frac{z_T \sinh(\omega(t - t_0)) + z_0 \sinh(\omega(T - t))}{\sinh(\omega(T - t_0))} \quad (14)$$

and the control is given by

$$a^*(t) = \omega \frac{z_T \cosh(\omega(t - T)) - z_0 \cosh(\omega(T - t))}{\sinh(\omega(T - t_0))} \quad (15)$$

where  $\omega^2 = -\frac{u_0}{K}$ .

If  $u_0 > 0$ , the cost function is given by (16), the optimal trajectory is

$$z^*(t) = \frac{z_T \sin(\omega(t - t_0)) + z_0 \sin(\omega(T - t))}{\sin(\omega(T - t_0))} \quad (17)$$

and the control is given by

$$a^*(t) = \omega \frac{z_T \cos(\omega(t - t_0)) - z_0 \cos(\omega(T - t))}{\sin(\omega(T - t_0))} \quad (18)$$

where  $\omega^2 = \frac{u_0}{K}$ .

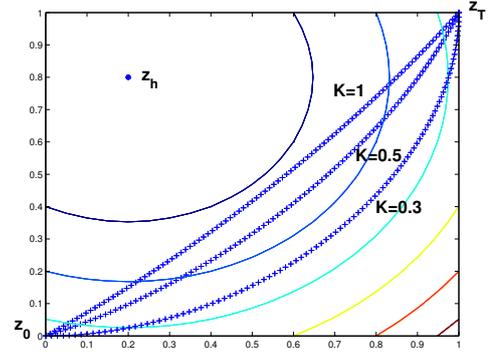
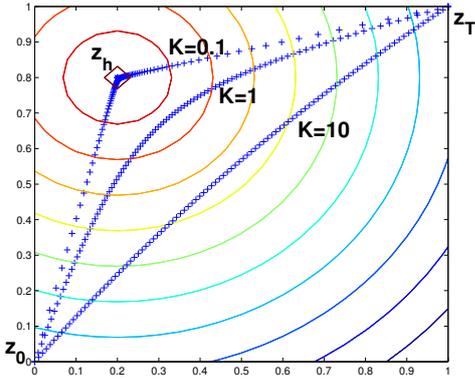
**Corollary 1.** *If the user traffic intensity is of the form  $u(t, z) = \frac{1}{2}u_0||z||^2 + u_0z \cdot e + u_1$  with  $u_0 \in \mathbb{R}$ ,  $u_1 \in \mathbb{R}$  and  $e \in \mathbb{R}^2$ , then define  $\tilde{z} = z + e$ ,  $\tilde{z}_0 = z_0 + e$ ,  $\tilde{z}_T = z_T + e$  and trajectories given in Theorem 1 are valid by replacing  $z$ ,  $z_0$ ,  $z_T$  by  $\tilde{z}$ ,  $\tilde{z}_0$ ,  $\tilde{z}_T$ , respectively. The cost function becomes:  $S(t_0, z_0, T, z_T) = \frac{1}{\alpha} [z \cdot p]_{t_0}^T + J(z_T) - u_1(T - t_0)$ .*

**Corollary 2.** *If the user traffic intensity is of the form  $u(t, z) = \sum_i u_i ||z - z_i||^2$  with  $\sum_i u_i \neq 0$ , then  $u(t, z) = (\sum_i u_i) ||z - z_b||^2 + \sum_i u_i ||z_i - z_b||^2$  with  $z_b = \frac{\sum_i u_i z_i}{\sum_i u_i}$ . Define  $\tilde{z} = z + z_b$ ,  $\tilde{z}_0 = z_0 + z_b$ ,  $\tilde{z}_T = z_T + z_b$ ,  $\tilde{u}_0 = \sum_i u_i$  and trajectories given in Theorem 1 are valid by replacing  $z$ ,  $z_0$ ,  $z_T$ ,  $u_0$  by  $\tilde{z}$ ,  $\tilde{z}_0$ ,  $\tilde{z}_T$ ,  $\tilde{u}_0$  respectively.*

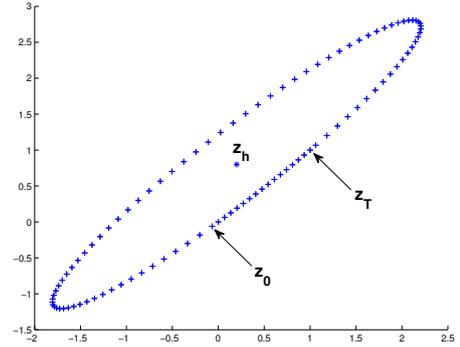
<sup>2</sup>We leave for further work the way to approximate a realistic traffic intensity map by a set of regions with intensities of quadratic form.

$$S(t_0, z_0, T, z_T) = \frac{K\omega}{2 \sinh \omega(T - t_0)} ( (|z_0|^2 + |z_T|^2) \cosh \omega(T - t_0) - 2z_0 \cdot z_T ) + J(z_T) \quad (13)$$

$$S(t_0, z_0, T, z_T) = \frac{K\omega}{2 \sin \omega(T - t_0)} ( (|z_0|^2 + |z_T|^2) \cos \omega(T - t_0) - 2z_0 \cdot z_T ) + J(z_T) \quad (16)$$



(a)  $T$  is smaller than the ellipse period.



(b)  $T$  is larger than the ellipse period.

Fig. 3: Traffic hole ( $u_0 > 0$ ).

Fig. 2: Traffic hot spot ( $u_0 < 0$ ). Circles are iso-traffic levels.

The system is thus equivalent to the one assumed in Theorem 1 by changing the origin of the locations to the barycentre  $z_b$  of the  $z_i$ .

2) *Traffic Hot Spot, Traffic Hole*: We assume that there is a hot spot or a traffic hole located in  $z_h$  and that the traffic intensity is of the form  $u(t, z) = \frac{1}{2}u_0||z - z_h||^2 + u_1 = \frac{1}{2}u_0||z||^2 - u_0z \cdot z_h + \frac{1}{2}u_0||z_h||^2 + u_1$ . We can apply Corollary 1 with  $e = -z_h$ . Figure 2 shows optimal trajectories when  $z_h$  is a hot spot, i.e., for  $u_0 < 0$ , and different values of  $K$ . The starting point is  $z_0$  and the destination is  $z_T$ . When  $K$  increases, the velocity cost increases and the trajectories tend to the straight line between  $z_0$  and  $z_T$ , which minimizes the speed. When  $K$  is small, the UAV can go fast to  $z_h$ , reduces its speed in the vicinity of the hot spot and then goes fast to the destination. Figure 3 shows optimal trajectories when  $z_h$  is a traffic hole, i.e., for  $u_0 > 0$ . In Figure 3a,  $T$  is smaller than the period of the ellipse, i.e.,  $\frac{2\pi}{\omega} > T$ . When  $K$  decreases, the UAV can spend more time in the areas of higher traffic intensity. In Figure 3b,  $T$  is larger than the period. In this case, the trajectory follows one period of the ellipse whose equation is given by (17) plus a part of the same ellipse from  $z_0$  to  $z_T$ .

### B. Multi-Phase Trajectory Characterization

We now consider a traffic intensity (or potential) consisting in two quadratic functions separated by an interface  $\mathcal{I}$  of equal potentials delimiting two regions 1 and 2. The interface is defined by an equation  $f(z) = C$ , where  $C$  is a constant and  $f$  is a differentiable function. We assume that the optimal trajectory crosses only once the interface at position  $\xi$  at  $\tau$ .

**Theorem 2.** *The location and time  $(\xi, \tau)$  of interface crossing are characterized by the following equations*

$$H_1(\xi(\tau), p^*(\tau^-)) - H_2(\xi(\tau), p^*(\tau^+)) = 0 \quad (19)$$

$$p^*(\tau^-) - p^*(\tau^+) - \mu \nabla_z f(\xi) = 0 \quad (20)$$

$$f(\xi) = C \quad (21)$$

for some Lagrange multiplier  $\mu \in \mathbb{R}$ , where we recall that  $p^*$  is defined with respect to the optimal trajectory between  $(t_0, z_0)$  and  $(T, z_T)$ , and where  $p^*(\tau^-) = \lim_{s \rightarrow \tau, s < \tau} p^*(s)$  and  $p^*(\tau^+) = \lim_{s \rightarrow \tau, s > \tau} p^*(s)$ .

Equation (19) expresses the fact the energy is conserved when crossing the interface. One can show that actually the energy is conserved along the whole trajectory. Equation (20) is related to the conservation of the tangential component of the impulsions at the interface. Equation (21) is the interface equation at  $\xi$ . One can show that under the assumption of equal potential on the interface, the kinetic energy, the

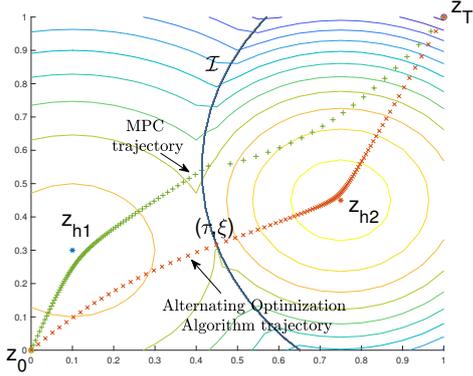


Fig. 4: MPC trajectory and Alternating Optimization Algorithm trajectory with two hot spots.

impulsion, and the velocity vector are conserved across the interface.

## V. ALGORITHMS

### A. An Online Algorithm: MPC

In this section, we first present an online algorithm based on MPC [18] (we omit the pseudo-code for space reasons). In a traffic intensity landscape made of multiple phases, the idea is to assume at every  $t$  that the current phase won't change from  $t$  to  $T$ . Using this assumption, we compute the optimal trajectory as in the single phase case and take the next decision based on this trajectory. This algorithm has the advantage of being online, of low complexity and can be used in multiphase time-dependent traffic maps. We have however no guarantee of optimality.

### B. An Alternating Optimization Algorithm

We now study a time-independent bi-phase scenario, in which a trajectory from  $z_0$  to  $z_T$  crosses the interface at time  $\tau$  and location  $\xi$ . We present an Alternating Optimization Algorithm (Algorithm 1) that provides a stationary trajectory in the sense of Theorem 2. The algorithm consists in alternatively optimizing  $\tau$  (steps 9-17) and  $\xi$  (steps 18-26) by using the results of Theorem 2. For every fixed  $\tau$  and  $\xi$ , the current trajectory is the concatenation of the optimal trajectory between  $(t_0, z_0)$  and  $(\tau, \xi)$  and the optimal trajectory between  $(\tau, \xi)$  and  $(T, z_T)$  (step 27). Every iteration of the algorithm only requires the evaluation of two Hamiltonians or the computation of a point  $B$ , see (23), and its projection on the interface. Therefore the complexity of an iteration is  $O(1)$ . In simulations, MPC is used to produce an initial trajectory.

1) *Procedure for seeking an optimal  $\tau$  given a fixed  $\xi$ :* We use the result of Theorem 2. As shown in its proof [17], the gradient of  $S$  with respect to  $\tau$  is given by  $H_2(\xi, p^*(\tau^+)) - H_1(\xi, p^*(\tau^-))$ . We can thus compute the Hamiltonians in every region by differentiating the cost function (13) with respect to the final time in region 1 (see (9)) and with respect to the initial time in region 2 (see (8)). We then update  $\tau$  by using a simple gradient descent scheme in step 11.

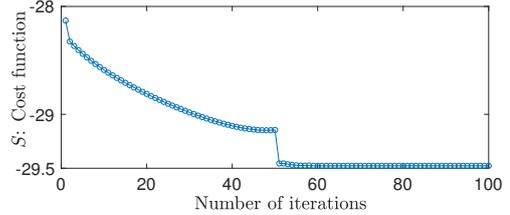


Fig. 5: Cost function along the iterations of the Alternating Optimization Algorithm trajectory.

2) *Procedure for seeking an optimal  $\xi$  given a fixed  $\tau$ :* From Hamilton-Jacobi, the gradient of the total cost function with respect to  $\xi$  is  $p^*(\tau^-) - p^*(\tau^+)$  (see proof of Theorem 2 in [17]). Since in the Newtonian framework the impulsion is proportional to the control variable  $a$  (see (12)) and since in a quadratic model the velocity vector is, at any time a linear combination of *centered* initial and final positions (15), this gradient appears to be an *affine* function of  $\xi$  which reads

$$\nabla_{z_T} S_1(t_0, z_0, \tau, \xi) + \nabla_{z_0} S_2(\tau, \xi, T, z_T) = Kh (\xi - B)$$

Scalar Hessian  $h$  and position  $B$ , where the spatial gradient cancels *i.e.*,  $p^*(\tau^-) = p^*(\tau^+)$  at fixed  $\tau$  are given by:

$$h = \omega_1 \coth(\omega_1(\tau - t_0)) + \omega_2 \coth(\omega_2(T - \tau)) \quad (22)$$

$$B = \frac{1}{h} \left[ \omega_1 z_{h1} \coth(\omega_1(\tau - t_0)) + \omega_2 z_{h2} \coth(\omega_2(T - \tau)) + \frac{\omega_1 (z_0 - z_{h1})}{\sinh(\omega_1(\tau - t_0))} + \frac{\omega_2 (z_T - z_{h2})}{\sinh(\omega_2(T - \tau))} \right] \quad (23)$$

The equation involving the Lagrange multiplier (20) now reads

$$K h (\xi - B) - \mu \nabla_{\xi} f(\xi) = 0 \quad (24)$$

and shows that the optimal location  $\xi^*$  is the *orthogonal projection* of  $B$  on the interface  $\mathcal{I}$ . This projection is performed in steps 19-20 of the algorithm.

Figure 4 shows the MPC trajectory and the trajectory obtained from Algorithm 1 after 60 iterations in a bi-phase landscape. The traffic intensity is shown in three dimensions in Figure 1: It is a bi-phase landscape made of two hot-spots, where the peak of traffic in  $z_{h2}$  is higher than in  $z_{h1}$ . The Alternating Optimization Algorithm has gradually moved the interface crossing time and location in order to spend more time in the second hot-spot and to go closer to  $z_{h2}$ . Figure 5 shows how the cost function has decreased along the iterations and thus how our algorithm has improved over the MPC solution. From iterations 1 to 45,  $\tau$  has been gradually updated; at iteration 46,  $\xi$  is updated once;  $\xi$  is again updated once at iteration 59.

## VI. CONCLUSION

In this paper, we have proposed a Lagrangian approach to solve the UAV base station optimal trajectory problem. When the traffic intensity exhibits a single phase, closed-form expressions for the trajectory and speed are given. When the traffic intensity exhibits multiple phases, we characterize the crossing time and location at the interface. In a first

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**Algorithm 1** Alternating Optimization Algorithm

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1: Input:  $t_0, T, z_0, z_T, z_{h1}, z_{h2}, u_{01}, u_{02}, \omega_1, \omega_2, u_{11},$   
    $u_{12}$ , an initial trajectory  $z(t)$ , the initial crossing time  
   and position  $(\tau, \xi) \in [\tau; T] \times \mathcal{I}$ ,  $\delta\tau > 0$ ,  $\epsilon_\tau > 0$ ,  $\epsilon_\xi > 0$ ,  
    $\epsilon_S > 0$ .  
2: Output:  $(\tau, \xi) \in [\tau; T] \times \mathcal{I}$  such that the conditions of  
   Theorem 2  
3:  $\tau' \leftarrow \tau$ ;  $\xi' \leftarrow \xi$   
4: timenotfound  $\leftarrow 1$ ; positionnotfound  $\leftarrow 0$   
5:  $\{z(t)\}_{t_0 \leq t \leq T} \leftarrow$  an initial feasible trajectory, e.g. from  
   MPC  
6: Compute  $S$  along  $\{z(t)\}_{t_0 \leq t \leq T}$   
7: do  
8:    $S' \leftarrow S$   
9:   if timenotfound then  
10:    Compute  $H_1$  and  $H_2$  at  $(\tau, \xi)$  according to (8-9)  
11:     $\tau \leftarrow \tau + \text{sign}(H_1 - H_2)\delta\tau$   
12:    if  $|\tau' - \tau| < \epsilon_\tau$  then  
13:      timenotfound  $\leftarrow 0$   
14:      positionnotfound  $\leftarrow 1$   
15:    end if  
16:     $\tau' \leftarrow \tau$   
17:  end if  
18:  if positionnotfound then  
19:    Compute  $B$  according to (23)  
20:     $\xi \leftarrow \text{proj}_{\mathcal{I}}(B)$ , see (24)  
21:    if  $\|\xi' - \xi\| < \epsilon_\xi$  then  
22:      timenotfound  $\leftarrow 1$   
23:      positionnotfound  $\leftarrow 0$   
24:    end if  
25:     $\xi' \leftarrow \xi$   
26:  end if  
27:   $\{z(t)\}_{t_0 \leq t \leq T} \leftarrow \text{OPTTRAJ}(z_{h1}, u_{01}, u_{11}, \omega_1, z_0, t_0,$   
    $\xi, \tau) \cup \text{OPTTRAJ}(z_{h2}, u_{02}, u_{12}, \omega_2, \xi, \tau, z_T, T)$   
   (OPTTRAJ provides optimal trajectory using (14),(17))  
28:  Compute  $S$  for  $\{z(t)\}_{t_0 \leq t \leq T}$  according to (13)  
29: while  $|S' - S| > \epsilon_S$ 
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approach, we propose an online algorithm based on MPC for multi-phase and time-dependent traffic intensity, which allows to take into account the impact of each phase. We then propose an offline Alternating Optimization Algorithm for bi-phase time-independent traffic intensities that provides a stationary trajectory with respect to the crossing time and location on the interface and fulfills the necessary conditions of optimality. Numerical results show that we improve the trajectory obtained with MPC.

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